

# Typed Functional Programming In OCaml

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# Introducing Myself

- Full-time researcher at INRIA, programming languages and distributed systems
- 2001 : PhD on **JoCaml**, a DSL for concurrency, distribution and mobility
- 2002 : **MLdonkey**, first multi-protocol peer-to-peer client (edonkey, gnutella, bittorrent, etc.)
- 2007 : **MNPLight**, first iPhone application able to install mp3s on a jailbroken iPhone 1
  - all in the **OCaml programming language**

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- 2011 : **OCamlPro**, a company to support the use of OCaml in industrial projects

# A Poll !

How many of you have some experience of:

- Lisp or Scheme ?
- F# or Scala ?
- Haskell ?
- OCaml ?

# What is OCaml ?

- A **General-purpose** Programming Language **developed** for about **30 years** at INRIA
  - **Used** *from the beginning* to **develop** many applications at INRIA :
    - **Coq proof assistant**, Zenon, Alt-ergo
    - Hevea (LaTeX → Html)
    - spamoracle (bayesian spam filter)
    - synDEX (scheduler for embedded systems)
    - Coccinelle (Linux Kernel bug checker)
- **OCaml is definitively not a lab toy !**

# OCaml, as a FP language (1)

- What is Functional Programming ?
  - A way of programming, closer to mathematics
    - make it easier to implement complex algorithms
    - make it possible to reason about the correctness of implementations
  - Usual features of FP languages:
    - immutable variables, immutable values
    - functions as values
    - use of (tail) recursion instead of loops
    - strong type-checking

# OCaml, as a FP language (2)

- Where is OCaml among FP languages:
  - **Hybrid FP languages:** Scala, F#, Clojure, etc.
    - FP extensions, "a taste of FP"... but tainted
  - **Untyped FP languages:** Lisp, Scheme, Erlang, etc.
    - FP, lack the power of strong type systems
  - **Pragmatic FP languages:** OCaml, SML
    - add other styles over FP, best of both worlds ?
  - **Pure FP languages:** Haskell
    - closer to maths, but hard to program with
  - **Proof languages:** Coq, Isabelle, etc.
    - write a math proof, generate code from it

# OCaml in the Industry

OCaml was designed at the beginning for formal methods applications: compilers, verifiers, provers...

- Microsoft : SLAM driver verifier
- Esterel Technologies : Scade KCG Compiler (scade-to-C, qualified level A DO-187B)
- AbsInt : Astree no-RTE checker
- EADS : Penjili, C code checker
- Dassault Systemes : Lucid/Esterel Compiler
- Airbus/Atos Origin: Toaster C style-checker

# Success Stories: Citrix

- 2002: Cambridge University releases Xen
  - need a program to control Xen in VM0
- 2004: 30 developers, C, Python et Ruby...
- 2006: many m\$ spent, yet, no product...
- 2006: new team of 4 OCaml devs, hired to write the doc, start a prototype in OCaml
- 2007: product available in OCaml,  
XenSource sold 500m\$ to Citrix
- 2011: Citrix holds 15% of the virtualisation market (Amazon EC2 for example)

# Success Stories: Jane Street

- 2000: Jane Street starts high-frequency trading in Excel + Visual Basic, too unreliable
- 2003: begin conversion from VB to C#
- 2003: one intern starts writing OCaml code
- 2005: management decides to try OCaml for key trading systems
- 2006: half of the system already in OCaml
- 2012: 10 billion\$ per day of automatic trading, everything in OCaml with 100+ OCaml devs

# OCaml is a multi-paradigm language

- **Functional** (functions are values, tail recursion)
- **Modular** (interfaces, functors and first-class modules)
- **Imperative** (mutable values, loops, exceptions)
- **Object-oriented** (objects and classes)
- **Statically and Strongly Typed**
- Execution is **strict** by default, **lazy** on demand
  - **Strict** = computation done where it is written
  - **Lazy** = computation delayed until useful

# OCaml Implementation

- **Native-code compiler** for x86/amd64, arm,...
- **Bytecode compiler, interpreter (REPL) and debugger** for fast development loop
- **Efficient incremental** garbage collector  
with **compaction**
- **Compact** uniform data **representation**
- Small but efficient **standard library**
- **FFI bindings** with many C libraries (databases, crypto, GUIs, etc.)

# Performances ?

- Strong Typing → No runtime checks !
- Highly optimised GC for short lifetime values
- Native-code compiler
  - with few but efficient optimisations  
(constant folding, inlining, register coalescing)
- Strict execution
  - expectable performance
  - close to non-optimized C speed  
(about 15% slower)
  - easy to optimise manually

# A Taste of OCaml

- **Warning:**

OCaml has a weird syntax

- Difficult to learn at the beginning... :-)
- Makes programs easier to read on the long term :-)

# Basic Values

## Simple values

```
let str = "Hello world"
let four = 2 * 2
let pi2 = 3.14 *. 2.    (* No operator overloading ! *)
let list = [ 1 ; 2 ; 3 ; 4 ; 5 ]
let list = 1 :: 2 :: 3 :: 4 :: 5 :: []
let tuple = (x, y, z)
let array = [| ('a', 97); ('b', 98); ('c', 99) |]
let record = { x = 1; y = 12 }
```

There are no NULL pointer in OCaml, all values **must be initialized !**

# Calling Functions

Functions arguments are **curried**:

```
let add (x,y) = x + y      (* one argument ! *)  
let add x y = x + y      (* two arguments ! *)  
let three = add (add 1 1) 1
```

Functions can be **partially applied**:

```
let add_one = add 1      (* val add_one : int → int *)  
  
let list = [ 1;2;3;4;5 ]  
let list_plus_two = List.map ( add 2 ) list  
                      (* → [ 3; 4; 5; 6; 7 ] *)
```

# Recursive Functions

Recursivity is intuitive to work on lists and trees

```
let rec fold_left f acc list =  
  match list with  
    [] → acc  
  | head :: tail →  
    fold_left f (f acc head) tail  
  (* fold_left f x [a;b;c] ↔ *)  
  (* f (f (f x a) b) c *)
```

```
let sum_list = List.fold_left add 0  
  (* val sum_list : int list → int *)  
let mul x y = x *. y  
let mul_list = List.fold_left mul 1.  
  (* val mul_list : float list → float *)
```

# Imperative Style

Side-effects, loops (while, for) and exceptions

```
let read_lines filename =  
  let ic = open_in filename in  
  let lines = ref [ ] in  
  try (* ref: mutable value *)  
    while true do  
      lines := input_line ic :: !lines  
      (* := modifies, ! for extraction *)  
    done; assert false  
  with End_of_file → (* exceptions are cost-free ! *)  
    close_in ic; List.rev !lines
```

# FP Style

```
let (>) x f = f x                (* x |> f ↔ f x *)
let normal_users = "/etc/passwd"
  |> read_lines
  |> List.map (Str.split_delim (Str.regexp ":"))
  |> List.map (fun list → match list with
    | login :: _passwd :: uid :: _ ->
      (login, int_of_string uid)
    | _ -> assert false )
  |> List.filter (fun (_, uid) -> uid >= 1000)
  |> List.map fst
```

# Closures

Variable bindings last for ever in functions...

```
let new_counter () =  
  let x = ref 0 in fun () → x := !x + 1; !x  
  (* new_counter : unit → (unit → int) *)  
let new_foo = new_counter ()  
let foo_id1 = new_foo () (* → 1 *)  
let new_bar = new_counter ()  
let bar_id1 = new_bar() (* → 1 *)  
let bar_id2 = new_bar() (* → 2 *)  
let bar_id3 = new_bar() (* → 3 *)  
let foo_id2 = new_foo() (* → 2 *)
```

# Where are type annotations ?

- OCaml is a **statically typed** language with one of the **most expressive type-systems** (variants, records, optional args, GADTs, polymorphic variants, objects, classes, etc.)
- Compiler is supposed to verify types !
  - but I didn't see any type annotations ?
- In OCaml, types are **automatically inferred**:
  - You don't need to write them
  - The compiler will guess them, and complain if they don't match what is expected

# Type-inference

```
(* val read_lines: string → string list *)
let read_lines filename =
  let ic = open_in filename in
    (* filename : string & ic : in_channel *)
  let lines = ref [ ] in
    (* lines: '_a list ref *)
  try
    while true do
      lines := input_line ic :: !lines
      (* lines: string list ref *)
    done; assert false
  with End_of_file -> close_in ic;
  List.rev !lines
```

# Polymorphic Functions

Our function on list works on any list !

```
let rec fold_left f acc list = match list with
  [] → acc
  | head :: tail →
    fold_left f (f acc head) tail
    (* ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a *)
```

```
let sum_list = List.fold_left add 0
    (* val sum_list : int list → int *)
let mul x y = x *. y
let mul_list = List.fold_left mul 1.
    (* val mul_list : float list → float *)
```

# Defining New Types

```
type expression =  
  { exp = expression_desc;  
    loc : Location.t; }  
and expression_desc =  
  | Num of int  
  | Var of string  
  | Let of string * expression * expression  
  | Binop of operator * expression * expression  
and operator = Plus | Minus | Times | Div
```

**Algebraic Data Types** avoid accessing the wrong arguments of an enum selector.

# Pattern-Matching

```
let rec eval env v = match v.desc with
| Num i -> i
| Var x -> List.assoc x env
| Let (x, e1, body) -> let val_x = eval env e1 in
    eval ((x, val_x) :: env) body
| Binop (Plus, e1, e2) ->
    (eval env e1) + (eval env e2)
| Binop (Minus, e1, e2) ->
    (eval env e1) - (eval env e2)
```

It is possible to match **deep and complex patterns**, that are always compiled in the **optimal number of runtime tests**.

# More Checks

```
let rec eval env v = match v.desc with
| Num i -> i
| Var x -> List.assoc x env
| Let (x, e1, body) -> let val_x = eval env e1 in
    eval ((x, val_x) :: env) body
| Binop (Plus, e1, e2) ->
    (eval env e1) + (eval env e2)
| Binop (Minus, e1, e2) ->
    (eval env e1) - (eval env e2)
```

Warning 8: **this pattern-matching is not exhaustive.**  
Here is an **example** of a value that is not matched:  
**Binop (Times | Div, \_, \_)**

# Simple Networking

```
(* start_server : int → (int → unit) → unit *)  
let start_server port handle_connection =  
  let server = Unix.socket PF_INET SOCK_STREAM 0 in  
  Unix.setsockopt server SO_REUSEADDR true;  
  Unix.bind server ADDR_INET (inet_addr_any, port);  
  Unix.listen server 3;  
  while true do  
    let (client, addr) = Unix.accept server in  
    ignore (Thread.create handle_connection client)  
  done
```

# Modules and Interfaces

Interface file: server.mli

```
val start_server : int → (int → unit) → unit
val read_lines : string → string list
...
```

Implementation file: server.ml

```
open Unix
let read_lines filename = ...
let start_server port handle_connection = ...
```

The compiler checks the **consistency of all compiled files** in the **whole project**: the compiler is often used as a **refactoring assistant** !

# OCaml Ecosystem

- OPAM, a **source package manager** to install OCaml and its open-source contributions
  - <http://opam.ocamlpro.com/>
- Js\_of\_ocaml, a powerful **OCaml-to-JavaScript optimizing compiler**, to run OCaml typed-checked applications in the browser
  - [http://oscigen.org/js\\_of\\_ocaml/](http://oscigen.org/js_of_ocaml/)
- Mirage: **bare-metal applications for Xen** in OCaml, speed and security in a **Cloud OS** !
  - <http://openmirage.org/>

# Formal Methods

- Use of Mathematics in the design of Hardware/Software applications
- **Strong type-checking** with OCaml
- **Abstract Interpretation:**
  - Astree, no runtime error in Airbus C code  
<http://www.absint.com/astree/>
- **Verification of formal specifications:**
  - Frama-C: used by Airbus on critical boot code  
<http://frama-c.com/>

# Formal Methods

- **Mecanized Proof of an Algorithm and automatic code generation:**
  - CompCert, a full C compiler, proved within the Coq proof assistant
  - <http://compcert.inria.fr/>
  - <http://coq.inria.fr/>

# Discussion

- Questions ?
- OCaml:
  - Web: <http://www.ocaml.org/>
  - Try it online: <http://try.ocamlpro.com/>
  - Install: <http://opam.ocamlpro.com/>